Scenario reduction and scenario tree construction for power management problems

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1. Electricity portfolio management

We consider a power utility owning a hydro-thermal generation system and producing and trading electric power.

Objective: Maximization of (expected) revenue

Decisions: Mixed-integer (large scale)

System and trading constraints: Capacity, reservoir, operational, load, reserve (and risk) constraints

Stochastic data processes: Electrical load, fuel and electricity prices, inflows



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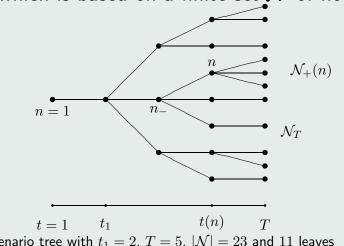
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Data process approximation by scenario trees 1.1.

The data process $\xi = \{\xi_t\}_{t=1}^T$ is approximated by a process forming a scenario tree which is based on a finite set $\mathcal N$ of nodes.



Scenario tree with $t_1=2$, T=5, $|\mathcal{N}|=23$ and 11 leaves

The root node n=1 stands for period t=1. Every other node n has a unique predecessor n_- and a set $\mathcal{N}_+(n)$ of successors. Let path(n) be the set $\{1, \ldots, n_-, n\}$ of nodes from the root to node $n, t(n) := |\operatorname{path}(n)|$ and $\mathcal{N}_T := \{n \in \mathcal{N} : \mathcal{N}_+(n) = \emptyset\}$ the set of leaves. A scenario corresponds to path(n) for some $n \in \mathcal{N}_T$. With the given scenario probabilities $\{\pi_n\}_{n \in \mathcal{N}_T}$, we define recursively node probabilities $\pi_n := \sum_{n_+ \in \mathcal{N}_+(n)} \pi_{n_+}, \ n \in \mathcal{N}$.



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1.2. Stochastic power management model

Stochastic process: $\{\xi_t = (d_t, r_t, \gamma_t, \alpha_t, \beta_t, \zeta_t)\}_{t=1}^T$ (electrical load, spinning reserve, inflows, (fuel or electricity) prices) given as a (multivariate) scenario tree Mixed-integer programming problem:

$$\begin{split} \min \sum_{n \in \mathcal{N}} \pi_n \sum_{i=1}^{I} [C_i^n(p_i^n, u_i^n) + S_i^n(u_i)] \quad \text{s.t.} \\ p_{it(n)}^{\min} u_i^n &\leq p_i^n \leq p_{it(n)}^{\max} u_i^n, \quad u_i^n \in \{0, 1\}, \quad n \in \mathcal{N}, \ i = 1 \colon I, \\ u_i^{n_{-\tau}} - u_i^{n_{-(\tau+1)}} &\leq u_i^n, \quad \tau = 1 \colon \bar{\tau}_i - 1, \ n \in \mathcal{N}, \ i = 1 \colon I, \\ u_i^{n_{-(\tau+1)}} - u_i^{n_{-\tau}} &\leq 1 - u_i^n, \quad \tau = 1 \colon \underline{\tau}_i - 1, \ n \in \mathcal{N}, \ i = 1 \colon I, \\ 0 &\leq v_j^n \leq v_{jt(n)}^{\max}, \ 0 \leq w_j^n \leq w_{jt(n)}^{\max}, \ 0 \leq l_j^n \leq l_{jt(n)}^{\max}, \ n \in \mathcal{N}, \ j = 1 \colon J, \end{split}$$

pdflatex

$$l_{j}^{n} = l_{j}^{n-} - v_{j}^{n} + \eta_{j} w_{j}^{n} + \gamma_{j}^{n}, \quad n \in \mathcal{N}, \ j = 1: J,$$

$$l_{j}^{0} = l_{j}^{\text{in}}, \quad l_{j}^{n} = l_{j}^{\text{end}}, \quad n \in \mathcal{N}_{T}, \ j = 1: J,$$

$$\sum_{i=1}^{I} p_{i}^{n} + \sum_{j=1}^{J} (v_{j}^{n} - w_{j}^{n}) \ge d^{n}, \quad n \in \mathcal{N},$$

$$\sum_{i=1}^{I} (u_{i}^{n} p_{it(n)}^{\text{max}} - p_{i}^{n}) \ge r^{n}, \quad n \in \mathcal{N}.$$

 C_i^n are fuel or trading costs and S_i^n start-up costs of unit i at node $n \in \mathcal{N}$:

$$C_i^n(p_i^n, u_i^n) := \max_{l=1,\dots,\bar{l}} \{\alpha_{il}^n p_i^n + \beta_{il}^n u_i^n\}$$

$$S_i^n(u_i) := \max_{\tau = 0, \dots, \tau_i^c} \zeta_{i\tau}^n(u_i^n - \sum_{\kappa = 1}^{\tau} u_i^{n_{-\kappa}})$$



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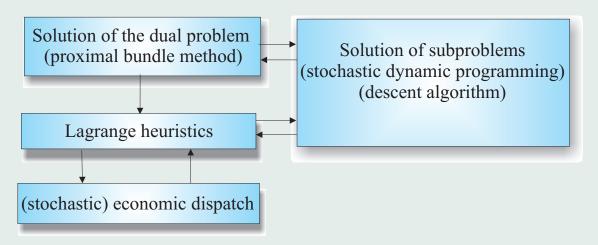
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1.3. Solving the stochastic power management model

$ \mathcal{N}_T $	$ \mathcal{N} $	va	riables	constraints	nonzeros
		binary	continuous		
1	168	4200	6652	13441	19657
20	1176	29400	45864	94100	137612
50	2478	61950	96642	198290	289976
100	4200	105000	163800	336100	491500

Dimension of the model for T=168, I=25 and J=7

- ⇒Primal approaches seem to be hopeless!
- ⇒Lagrangian relaxation of coupling constraints





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2. Generation of scenario trees

- (i) Development of a stochastic model for the data process ξ (parametric [e.g. time series model], nonparametric [e.g. resampling]) and generation of simulation scenarios;
- (ii) Construction of a scenario tree out of the stochastic model or of the simulation scenarios;
- (iii) optional scenario tree reduction.

Approaches for (ii):

- (1) Barycentric scenario trees (conditional expectations w.r.t. a decomposition of the support into simplices);
- (2) Fitting of trees with prescribed structure to given moments;
- (3) Conditional sampling by integration quadratures;
- (4) Clustering methods for bundling scenarios;
- (5) Scenario tree construction based on optimal approximations w.r.t. certain probability metrics.



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3. Distances of probability distributions

Let P denote the probability distribution of the stochastic data process $\{\xi_t\}_{t=1}^T$, where ξ_t has dimension r, i.e., P has support $\Xi \subseteq \mathbb{R}^{rT} = \mathbb{R}^s$.

The Kantorovich functional or transportation metric takes the following form

$$\mu_c(P,Q) := \inf\{ \int_{\Xi \times \Xi} c(\xi, \tilde{\xi}) \eta(d\xi, d\tilde{\xi}) : \pi_1 \eta = P, \ \pi_2 \eta = Q \},$$

where $c:\Xi\times\Xi\to I\!\! R$ is a certain cost function.

$$\underline{\mathsf{Example:}}\ c(\xi,\tilde{\xi}) := \max\{1, \|\xi\|^{p-1}, \|\tilde{\xi}\|^{p-1}\} \|\xi - \tilde{\xi}\| \quad (p \geq 1)$$

Approach:

Select a probability metric a function $c:\Xi\times\Xi\to I\!\!R$ such that the underlying stochastic optimization model is stable w.r.t. μ_c .

Given P and a tolerance $\varepsilon > 0$, determine a scenario tree such that its probability distribution P_{tr} has the property

$$\mu_c(P, P_{tr}) \leq \varepsilon$$
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Distances of discrete distributions

P: scenarios ξ_i with probabilities p_i , $i=1,\ldots,N$,

Q: scenarios ξ_i with probabilities q_i , $j = 1, \ldots, M$.

Then it holds

$$\mu_c(P,Q) = \sup\{\sum_{i=1}^{N} p_i u_i + \sum_{j=1}^{M} q_j v_j : u_i + v_j \le c(\xi_i, \tilde{\xi}_j) \ \forall i, j\}$$

$$= \inf \{ \sum_{i,j} \eta_{ij} c(\xi_i, \tilde{\xi_j}) : \eta_{ij} \ge 0, \sum_j \eta_{ij} = p_i, \sum_i \eta_{ij} = q_j \}$$

(linear transportation problem)

Optimal scenario generation:

Given P and $\varepsilon > 0$, determine $\bar{\xi}_j \in \Xi$, $\bar{q}_j \in [0,1]$, $j=1,\ldots,M$, such that it holds

$$\mu_c(P, \sum_{j=1}^M \bar{q}_j \delta_{\bar{\xi}_j}) \le \varepsilon.$$

(nonconvex facility location problem)



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4. Scenario Reduction

We consider the transportation metric μ_c and distributions P and Q as described above.

Theorem: (optimal reduction of a scenario set J)

$$D_J := \min_{q_j \ge 0, \sum_j q_j = 1} \mu_c(\sum_{i=1}^N p_i \delta_{\xi_i}, \sum_{j \notin J} q_j \delta_{\xi_j}) = \sum_{i \in J} p_i \min_{j \notin J} c(\omega_i, \omega_j)$$

The minimum is attained at $\bar{q}_j = p_j + \sum_{i \in J_j} p_i, \, \forall j \not\in J$, where

$$J_j := \{i \in J : j = j(i)\} \text{ and } j(i) \in \arg\min_{j \notin J} \vec{c}(\omega_i, \omega_j), \, \forall i \in J.$$

(optimal redistribution)

Optimal reduction of a scenario set with fixed cardinality:

$$\min\{D_J = \sum_{i \in J} p_i \min_{j \notin J} c(\xi_i, \xi_j) : J \subset \{1, ..., N\}, \#J = N - n\}$$

(combinatorial optimization problem of set-covering type)



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5. Fast reduction heuristics



Algorithm 1: (Simultaneous backward reduction)

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Step [0]: Sorting of $\{c(\xi_j, \xi_k) : \forall j\}, \forall k,$

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$$J^{[0]} := \emptyset.$$

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Step [i]: $l_i \in \arg\min_{l \notin J^{[i-1]}} \sum_{k \in J^{[i-1]} \cup \{l\}} p_k \min_{j \notin J^{[i-1]} \cup \{l\}} c(\xi_k, \xi_j).$

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$$J^{[i]} := J^{[i-1]} \cup \{l_i\}$$
.



Step [N-n+1]: Optimal redistribution.

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Algorithm 2: (Fast forward selection)

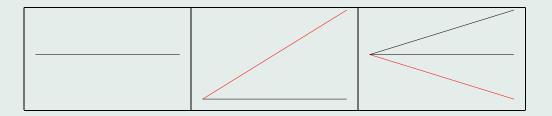
Step [0]: Compute $c(\xi_k, \xi_u), k, u = 1, ..., N,$

 $J^{[0]} := \{1, \dots, N\}.$

Step [i]: $u_i \in \arg\min_{u \in J^{[i-1]}} \sum_{k \in J^{[i-1]} \setminus \{u\}} p_k \min_{j \notin J^{[i-1]} \setminus \{u\}} c(\xi_k, \xi_j),$

$$J^{[i]} := J^{[i-1]} \setminus \{u_i\}.$$

Step [n+1]: Optimal redistribution.





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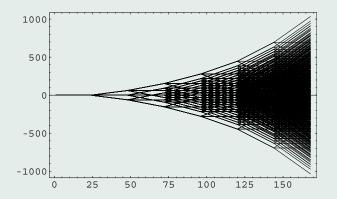


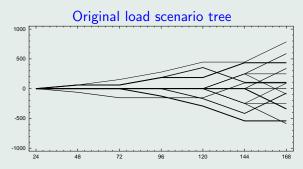
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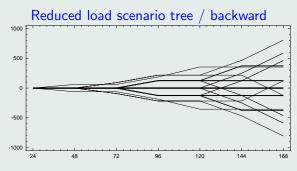
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Reduced load scenario tree / forward



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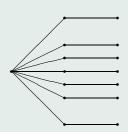
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6. Constructing scenario trees from data scenarios

Let a fan of data scenarios $\xi^i=(\xi^i_1,\ldots,\xi^i_T)$ with probabilities π^i , $i=1,\ldots,N$, be given, i.e., all scenarios coincide at the starting point t=1, i.e., $\xi^1_1=\ldots=\xi^N_1=:\xi^*_1$. Hence, it has the form



t=1 may be regarded as the root node of a scenario tree consisting of N branches (leaves, scenarios).

Now, P is the (discrete) probability distribution of ξ . Let c be adapted to the underlying stochastic program containing P.

We describe an **algorithm** that produces, for each $\varepsilon > 0$, a scenario tree with distribution P_{ε} , root node ξ_1^* , less nodes than P and

$$\mu_c(P, P_{\varepsilon}) < \varepsilon.$$



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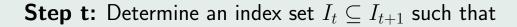
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Algorithm:

Let $\varepsilon_t > 0$, $t = 1, \ldots, T$, be given such that $\sum_{t=1}^T \varepsilon_t \leq \varepsilon$, set t := T, $I_{T+1} := \{1, \ldots, N\}$, $\pi^i_{T+1} := \pi^i$ and $P_{T+1} := P$. For $t = T, \ldots, 2$:



$$\mu_{c_t}(P_t, P_{t+1}) < \varepsilon_t \,,$$

where $\{\xi^i\}_{i\in I_t}$ is the support of P_t and c_t is defined by $c_t(\xi,\tilde{\xi}):=c((\xi_1,\ldots,\xi_t,0,\ldots,0),(\tilde{\xi}_1,\ldots,\tilde{\xi}_t,0,\ldots,0));$

(scenario reduction w.r.t. $\{1,\ldots,t\}$)

Step 1: Determine a probability measure P_{ε} such that its marginal distributions $P_{\varepsilon}\Pi_t^{-1}$ are $\delta_{\xi_1^*}$ for t=1 and

$$P_{\varepsilon}\Pi_t^{-1} = \sum_{i \in I_t} \pi_t^i \delta_{\xi_t^i} \quad \text{ and } \quad \pi_t^i := \pi_{t+1}^i + \sum_{j \in J_{t,i}} \pi_{t+1}^j \,,$$

where $J_{t,i} := \{j \in I_{t+1} \setminus I_t : i_t(j) = i\}, i_t(j) \in \arg\min_{i \in I_t} c_t(\xi^j, \xi^i)\}$ are the index sets according to the redistribution rule.



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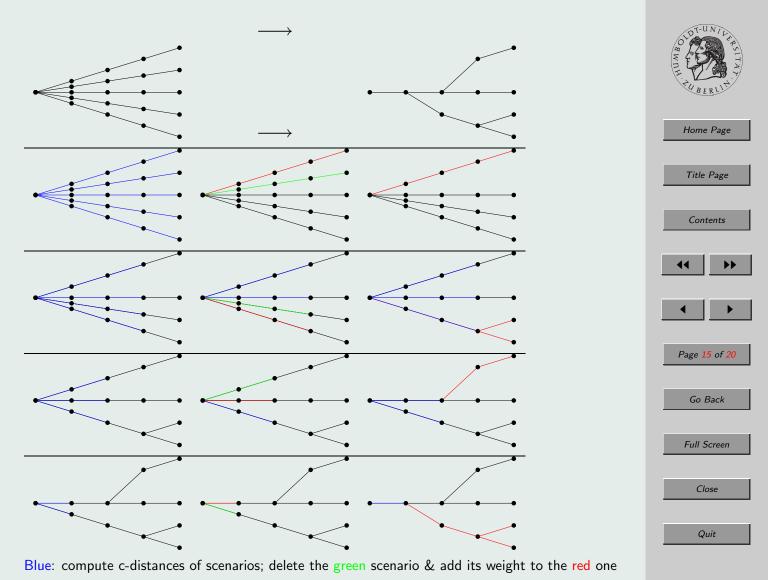
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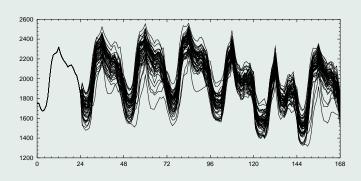
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Application:

- ξ is the multivariate data process having the components
- a) electrical load,
- b) electricity prices for baseload contracts (at EEX),
- c) electricity prices for peakload contracts (at EEX),
- d) electricity prices for individual hours (at EEX).

Data scenarios are obtained from a stochastic model calibrated to the historical load data of the German power utility VEAG and historical price data of the European Energy Exchange (EEX) at Frankfurt/Leipzig.



Scenario tree for the electrical load



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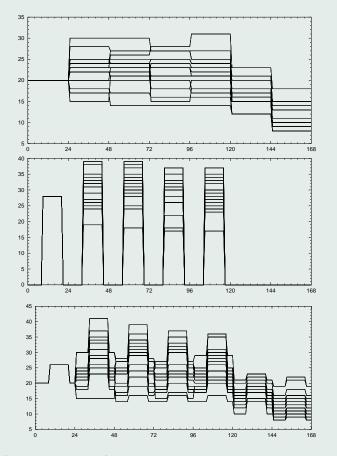
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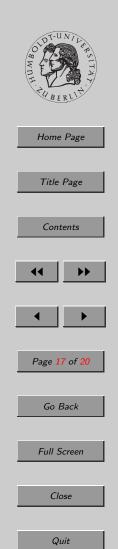
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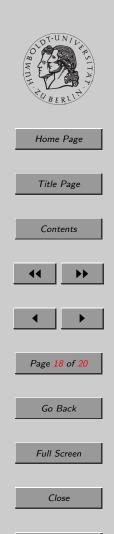


Scenario trees for prices of baseload contracts, peakload contracts and individual hours, respectively



7. GAMS/SCENRED

- GAMS/SCENRED introduced to GAMS Distribution 20.6 (May 2002)
- SCENRED is a collection of C++ routines for the optimal reduction of scenarios or scenario trees
- GAMS/SCENRED provides the link from GAMS programs to the scenario reduction algorithms. The reduced problems can then be solved by a deterministic optimization algorithm provided by GAMS.
- SCENRED contains three reduction algorithms:
 - FAST BACKWARD method
 - Mix of FAST BACKWARD/FORWARD methods
 - Mix of FAST BACKWARD/BACKWARD methods Automatic selection (best expected performance w.r.t. running time)



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Details: www.scenred.de, www.scenred.com

8. Numerical tests

We tested the link between the Lagrangian relaxation and the scenario tree construction algorithms.

- Portfolio management problem for 25 thermal generation units and 7 pumped-storage hydro units
- Time horizon: 1 week; Discretization: 1 hour
- Initial fan of 100 load scenarios simulated from a statistical model for the load process (combines a time series model for the daily mean load with regression models for the intra-day behaviour)







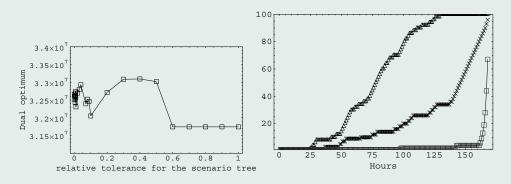






Dimension and solution time for the dual

$\varepsilon_{\mathrm{rel}}$	S	N	Variables		Nonzeros	time[s]
			binary	continuous		
0.6	1	168	4200	7728	44695	7.83
0.1	67	515	12875	23690	137459	17.09
0.05	81	901	22525	41446	240233	37.82
0.01	94	2660	66500	122360	708218	150.14
0.005	96	3811	95275	175306	1014398	291.65
0.001	100	9247	231175	425362	2460402	1176.38



Dual optimum and number of scenario bundles $|I_t|$ $(t=1,\ldots,T)$ for scenario trees with relative tolerance $\varepsilon_{\rm rel}=0.001$ (\triangle), 0.005 (\times), 0.01 (\square)



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